

Valid and Invalid Arguments

An argument is a sequence of statements such that

- all statements but the last are called *hypotheses*
- the final statement is called the *conclusion*.
- the symbol \therefore read “therefore” is usually placed just before the conclusion.

Example:

$$p \wedge \sim q \rightarrow r$$

$$p \vee q$$

$$q \rightarrow p$$

$$\therefore r$$

An argument is said to be **valid** if
- whenever all hypotheses are
true, the conclusion must be true.

Example of a valid argument (form)

$$p \wedge (q \vee r)$$

$$\sim q$$

$$\therefore p \wedge r$$

| p | q | r | $p \wedge (q \vee r)$ | $\sim q$ | $p \wedge r$ |
|---|---|---|-----------------------|----------|--------------|
| T | T | T | T | F | T |
| T | T | F | T | F | F |
| T | F | T | T | T | T |
| T | F | F | F | T | F |
| F | T | T | F | F | F |
| F | T | F | F | F | F |
| F | F | T | F | T | F |
| F | F | F | F | T | F |

An invalid argument

$$p \rightarrow q \vee \sim r$$

$$q \rightarrow p \vee r$$

$$\therefore p \rightarrow r$$

Invalid row

| p | q | r | $p \rightarrow q \vee \sim r$ | $q \rightarrow p \vee r$ | $p \rightarrow r$ |
|---|---|---|-------------------------------|--------------------------|-------------------|
| T | T | T | T | T | T |
| T | T | F | T | F | |
| T | F | T | F | T | |
| T | F | F | T | T | F |
| F | T | T | T | F | |
| F | T | F | T | F | |
| F | F | T | T | T | T |
| F | F | F | T | T | T |

Tautology

- is a statement (form) that is always true regardless of the truth values of the individual statement variables.

Examples:

- $p \vee \sim p$ (eg. the number n is either > 0 or ≤ 0)
- $p \wedge q \rightarrow p$
- $(p \rightarrow q \rightarrow r) \rightarrow (p \rightarrow r)$

We need to study tautologies because any valid argument is equivalent to a tautology.

In particular, every theorem we have proved is a tautology.

Example:

$$p \wedge (q \vee r)$$

$$\sim q$$

$$\therefore p \wedge r$$

is a valid argument,

$$[p \wedge (q \vee r)] \wedge [\sim q] \rightarrow [p \wedge r]$$

is a tautology.

In other words, an argument

$$H_1$$

$$H_2$$

...

$$H_n$$

\therefore Conclusion

is valid if and only if

$$H_1 \wedge H_2 \wedge \dots \wedge H_n \rightarrow \text{conclusion}$$

is a tautology.

Two most important valid argument forms

Modus Ponens : *means method of affirming*

$$p \rightarrow q$$

$$p$$

$$\therefore q$$

Example: If $n \geq 5$, then $n!$ is divisible by 10.

$$n = 7$$

$\therefore 7!$ is divisible by 10.

Modus Tollens : *means method of denying*

$$p \rightarrow q$$

$$\sim q$$

$$\therefore \sim p$$

Example: If n is odd, then n^2 is odd.

n^2 is even.

$\therefore n$ is even.

More valid forms

Conjunctive simplification: $p \wedge q$
 $\Leftrightarrow p$

Example: The function f is 1-to-1 and continuous.
 \Leftrightarrow The function f is 1-to-1.

Disjunctive addition: p
 $\Leftrightarrow p \vee q$

Example: The function f is increasing.
 \Leftrightarrow The function f is increasing or differentiable.

More valid forms

Conjunctive addition: p
 q
 $\text{∞ } p \text{ } \blacktriangle$

q
Example: n is an integer, n is positive.
 $\text{∞ } n$ is a positive integer.

Disjunctive syllogism: $p \text{ } \blacktriangledown \text{ } q$
" q
 $\text{∞ } p$

Example: The graph of this equation may be a circle or an ellipse.

The graph of this equation cannot be a circle.

$\text{∞ } \text{The graph must be an (true) ellipse.}$

Hypothetical syllogism:

$p \rightarrow q$
 $q \rightarrow r$
 $\therefore p \rightarrow r$

Proof by cases:

$p \vee \neg p$
 $p \rightarrow r$
 $\neg p \rightarrow r$
 $\therefore r$

Rule of contradiction:

$p \wedge \neg p$
 $\therefore p$

Example:

n is either odd or even.

If n is odd, then $n(n-1)$ is even.

If n is even, then $n(n-1)$ is even.

Therefore $n(n-1)$ is always even.

Example:

If $\sqrt{2}$ is not irrational, then there exists whole numbers a, b that are relatively prime and are both even.

A valid argument with a false conclusion.

The following argument is valid by modus ponens, but since its hypothesis is false, so is its conclusion.

If p is prime, then $2^p - 1$ is also prime (False).

11 is prime (True).

Therefore $2^{11} - 1$ is prime (False).

Actually, $2^{11} - 1 = 2047 = 23 \times 89$ is not prime.

Note: Any prime of the form $2^p - 1$ is called a Mersenne prime, the largest one up to date is $2^{6972593} - 1$ (discovered on 6-1-99)